

Hash Tables

Computing

How $O(1)$ lookup actually works, and when it stops being $O(1)$.

Every dictionary, map and set in every language is built this way.

1

Hash the key

A function turns the key into a large integer, deterministically.

2

Reduce to an index

Modulo the array size gives a slot.

3

Handle collisions

Two keys can land in one slot — chain them, or probe for the next free one.

4

Resize when full

Past a load factor around 0.7, allocate a bigger array and rehash everything.

WHEN IT DEGRADES

Many collisions $\rightarrow O(n)$ in the worst case

Resizing is $O(n)$, but amortised across inserts

A mutable key changes its hash and becomes unfindable

Never mutate an object after using it as a key. Its hash changes and the entry becomes unreachable.